Strong convergence analysis of iterative solvers for random operator equations

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Abstract

For the numerical solution of linear systems that arise from discretized partial differential equations, multigrid and domain decomposition methods are well established. Multigrid methods are known to have optimal complexity and domain decomposition methods are in particular useful for parallelization of the implemented algorithm. For random operator equations, the classical theory is not directly applicable, since condition numbers of system matrices may be close to degenerate due to non-uniform random input. As a main result, it is shown that iterative methods converge in the strong, i.e., L^p, sense if the random input satisfies certain integrability conditions. As a result, multigrid and domain decomposition methods are applicable in the case of elliptic partial differential equations with lognormal diffusion coefficients and converge strongly with deterministic bounds on the computational work which are essentially optimal. This enables the application of multilevel Monte Carlo methods with rigorous, deterministic bounds on the computational work.

1 Introduction

The efficient numerical approximation of solutions to stochastic equations is a key task in engineering applications that involve uncertainty. The present manuscript analyzes the applicability and strong convergence of well established iterative methods for operator equations.

Let $\mathcal{A}$ be a random, continuous linear operator from $\mathcal{V}$ to $\mathcal{V}^*$ on a probability space $(\Omega, \mathcal{F}, \mathbb{P})$ that is $\mathbb{P}$-almost surely ($\mathbb{P}$-a.s. for short) boundedly invertible, where $\mathcal{V}$ is a Hilbert space and $\mathcal{V}^*$ its dual space. Let its expectation $\mathbb{E}(\mathcal{A})$ be well defined. In the present paper we are interested in the numerical analysis of approximations of the solution u to the linear equation that for $\mathbb{P}$-almost every ($\mathbb{P}$-a.e. for short) $\omega \in \Omega$

$$\mathcal{A}(\omega)u(\omega) = f, \tag{1}$$

where $f \in \mathcal{V}^*$ is deterministic, in the strong sense by iterative methods such as multigrid and domain decomposition methods. This can be rewritten in variational form to find $u : \Omega \to \mathcal{V}$ such that for $\mathbb{P}$-a.e. ω

$$a_\omega(u(\omega), v) := {}_{\mathcal{V}^*}\langle \mathcal{A}(\omega)u(\omega), v\rangle_{\mathcal{V}} = {}_{\mathcal{V}^*}\langle f, v\rangle_{\mathcal{V}} \quad \forall v \in \mathcal{V}.$$

Let us assume that there are strictly positive random variables $\hat{a}$ and $\check{a}$ such that for $\mathbb{P}$-a.e. ω

$$\check{a}(\omega) \, {}_{\mathcal{V}^*}\langle \mathbb{E}(\mathcal{A})v, v\rangle_{\mathcal{V}} \le {}_{\mathcal{V}^*}\langle \mathcal{A}(\omega)v, v\rangle_{\mathcal{V}} \le \hat{a}(\omega) \, {}_{\mathcal{V}^*}\langle \mathbb{E}(\mathcal{A})v, v\rangle_{\mathcal{V}} \quad \forall v \in \mathcal{V}. \tag{2}$$

We will be particularly interested in the case that $\hat{a}$ and $\check{a}^{-1}$ are unbounded random variables. This is for example the case for elliptic partial differential equations (PDEs for short) with

lognormal coefficients. Thus, preconditioned finite dimensional discretizations suffer from random condition numbers that are unbounded and respectively interative methods contract with random contraction numbers, which may have realizations arbitrarily close to one with positive probability. At first sight one may overcome this with random interation numbers specified by a threshold of residual errors with the disadvantage of the occurance of large iteration numbers, when samples of the random contraction number are close to one. Also, bounds on the computational work for such strategy would be probabilistic. A main result of this paper is that deterministic iteration numbers exist that allow for strong convergence, i.e., convergence of certain iterative methods in the $L^q(\Omega; \mathcal{V})$-norm, $q \in [1, +\infty)$. This is possible due to tail bounds of the random contraction numbers, which for example are satisfied in the important case of elliptic PDEs with lognormal coefficients. In this case, deterministic, essentially optimal complexity bounds are implied for the solution of resulting random linear systems when well known multigrid or domain decomposition methods are applied. This enables also rigorous, deterministic complexity bounds for multilevel Monte Carlo (MLMC for short) approximations of mean fields. Assumptions on the computational cost of PDE solvers that were made in previous papers [12, 23, 18] to obtain complexity bounds of MLMC are pervaded by the theory presented in this manuscript. We will treat the case that $a_\omega(\cdot, \cdot)$ is symmetric for $\mathbb{P}$-a.e. ω. However, the presented theory can be extended to certain non-symmetric operators, see for example cp. [10, Section 11].

In Section 2, we will review iterative methods as they were formulated in [25] in order to discuss various multilevel method in a unified framework as well as the more classical approach, cp. [3, 28]. It will also be highlighted which parts in the framework and in the iterative methods are random. As a main result, we will develop integrability conditions on the random contraction numbers in Section 3 that result in sufficiently strong tail bounds in order to ensure strong convergence in the setting of multi and single-level discretizations of (1). The integrability conditions posed in Section 3 are analysed for several multilevel methods such as multigrid, the so called BPX preconditioner, cp. [8], and domain decomposition methods in Sections 4 and 5. An important application of the presented theory are lognormal diffusion problems, which are briefly reviewed in Section 6. In particular deterministic bounds on the computational work without assumptions on the PDE solver of MLMC are implied.

2 Iterative methods

In this section we recall well known iterative methods to approximate solutions to linear equations on a finite dimensional inner product space $(V, (\cdot, \cdot))$, where $\|\cdot\|$ denotes the norm that is induced by $(\cdot, \cdot)$. Let us consider the random linear equation that for $\mathbb{P}$-a.e. ω

$$A(\omega)u(\omega) = f, \tag{3}$$

where $A : V \to V$ is a random linear operator that is $\mathbb{P}$-a.s. symmetric and positive definite (SPD for short) with respect to $(\cdot, \cdot)$. Hence, (3) is $\mathbb{P}$-a.s. uniquely solvable. Note that we will often omit dependencies of random quantities on ω for notational convenience. Let us denote the bilinear form that is induced by A by $(\cdot, \cdot)_A$ and let $\lambda_{\max}(A)$, $\lambda_{\min}(A)$ denote the maximal and minimal eigenvalue of A. The condition number of A is denoted $\kappa(A)$ and $\rho(A)$ denotes the spectral radius. This notation will also be used for other linear operators that occur. Note that since A is random, $\lambda_{\max}(A)$, $\lambda_{\min}(A)$, $\kappa(A)$, and $\rho(A)$ are random variables. In particular, the review article [25] enables the discussion of multigrid and domain decomposition methods in an unified framework. These methods allow in some cases for optimal preconditioning or uniform contraction numbers. In this section we will mainly follow [25] and introduce abstract algorithms, which in later sections will be used as BPX or additive Schwarz preconditioner,

(un)symmetric multigrid and overlapping domain decomposition method. We will also highlight which of the occuring objects in this review section are random.

Since A is SPD $\mathbb{P}$-a.s., then the well known conjugent gradient (CG for short) method implies after $n \in \mathbb{N}$ iterations with initial guess U^0 the error bound that for $\mathbb{P}$-a.e. ω

$$\|u(\omega) - U^n(\omega)\|_{A(\omega)} \leq 2 \left(\frac{\sqrt{\kappa(A(\omega))} - 1}{\sqrt{\kappa(A(\omega))} + 1} \right)^n \|u(\omega) - U^0(\omega)\|_{A(\omega)}.$$

Since the random condition number $\kappa(A)$ may depend on the dimension of the linear space V, we consider the preconditioned linear system that for $\mathbb{P}$-a.e. ω

$$B(\omega)A(\omega)u(\omega) = B(\omega)f, \tag{4}$$

where the random linear operator B is chosen to be SPD with respect to $(\cdot, \cdot)$. The random operator B shall satisfy that $\mathbb{P}$-a.s. $\kappa(BA) \leq \kappa(A)$, which then accelerates the convergence of the CG method. The combination of preconditioning and CG will be referred to as preconditioned conjugent gradient (PCG for short) method.

Another method to be discussed is the linear iterative method of the form

$$U^{n+1}(\omega) = U^n(\omega) + B(\omega)(f - A(\omega)U^n(\omega)), \quad n \in \mathbb{N}_0, \tag{5}$$

for $\mathbb{P}$-a.e. ω, where B is a suitable random operator that is not nessarily symmetric and U^0 is given. Note that this linear iterative methods convergences $\mathbb{P}$-a.s. if $\| \operatorname{Id} - BA \|_A < 1$, $\mathbb{P}$-a.s. Alternatively, one could also introduce a relaxed version of B with relaxation parameter in $(0, 2/\rho(BA))$ to guarantee convergence with random contraction number $(\kappa(BA) - 1)/(\kappa(BA) + 1)$, cp. [25, Proposition 2.3]. Also we remark that generally the contraction number of the PCG method is smaller, cp. [25, Proposition 2.2], which is why one may say that PCG accelerates convergence.

Lemma 2.1 *Let $\tilde{A} : V \to V$ be a SPD operator with respect to $(\cdot, \cdot)$ and let $\tilde{B} : V \to V$ be a SPD preconditioner with respect to $(\cdot, \cdot)$ for $\tilde{A}$. If there exists positive random variables c_0, c_1 such that for $\mathbb{P}$-a.e. $\omega \in \Omega$*

$$c_0(\omega)(v, v)_{\tilde{A}} \leq (v, v)_{A(\omega)} \leq c_1(\omega)(v, v)_{\tilde{A}} \quad \forall v \in V,$$

then for $\mathbb{P}$-a.e. $\omega \in \Omega$

$$\kappa(\tilde{B}A(\omega)) \leq \kappa(\tilde{B}\tilde{A}) \frac{c_1(\omega)}{c_0(\omega)}.$$

Proof. Since $\tilde{B}\tilde{A}$ is SPD with respect to $(\cdot, \cdot)_{\tilde{A}}$, it holds that

$$\kappa(\tilde{B}\tilde{A}) = |\lambda_{\max}(\tilde{B}\tilde{A})/\lambda_{\min}(\tilde{B}\tilde{A})|.$$

According to [25, Lemma 2.1] this implies

$$|\lambda_{\max}(\tilde{B}\tilde{A})|^{-1}(v, v)_{\tilde{A}} \leq (\tilde{B}^{-1}v, v) \leq |\lambda_{\max}(\tilde{B}\tilde{A})|^{-1}(v, v)_{\tilde{A}}.$$

Then, the assumption of the lemma implies that

$$c_1^{-1}|\lambda_{\max}(\tilde{B}\tilde{A})|^{-1}(Av, v) \leq (\tilde{B}^{-1}v, v) \leq c_0^{-1}|\lambda_{\min}(\tilde{B}\tilde{A})|^{-1}(Av, v), \tag{6}$$

which implies the claim with another application of [25, Lemma 2.1]. $\qquad\square$

For $J \in \mathbb{N}$, let us assume a decomposition of V in subspaces $(V_j, j = 1, \ldots, J)$, i.e., it holds that $V_j \subset V$, $j = 1, \ldots, J$, and

$$V = \sum_{j=1}^{J} V_j. \tag{7}$$

We define the orthogonal projections $Q_j, P_j : V \rightarrow V_j$ for every $v \in V$ by

$$(Q_j v, w_j) := (v, w_j), \quad (AP_j v, w_j) := (Av, w_j) \quad \forall w_j \in V_j$$

and the operator $A_j : V_j \rightarrow V_j$ for every $v \in V_j$ by

$$(A_j v, w_j) := (Av, w_j) \quad \forall w_j \in V_j,$$

$j = 1, \ldots, J$. Consequently it holds for every $j = 1, \ldots, J$ that $A_j P_j = Q_j A$, which implies that if u is the random solution of (3), then $u_j := P_j u$ satifies for $\mathbb{P}$-a.e. ω

$$A_j(\omega) u_j(\omega) = f_j,$$

where $f_j := Q_j f$, $j = 1, \ldots, J$. Note that A_j, P_j are random whereas Q_j is deterministic, $j = 1, \ldots, J$. Let $R_j : V_j \rightarrow V_j$, $j = 1, \ldots, J$, be random SPD operators with respect to $(\cdot, \cdot)$, which shall approximate the inverse of A_j respectively. Thus,

$$B^a := \sum_{j=1}^{J} R_j Q_j \tag{8}$$

is also SPD, cp. [25, Lemma 3.1], and a candidate for a random preconditioner of A, where the corresponding method is known as *parallel subspace correction*.

Algorithm 1 Apply the PCG method to (4) with the random preconditioner B^a defined in (8).

A multiplicative algorithm, cp. [25, Algorithm 3.3], can also be defined.

Algorithm 2 Let U^0 be given and U^n be already obtained, then U^{n+1} is defined by

$$U^{n+j/J} := U^{n+(j-1)/J} + R_j Q_j (f - AU^{n+(j-1)/J}), \quad j = 1, \ldots, J.$$

We introduce $T_j := R_j A_j P_j = R_j Q_j A$ for every $j = 1, \ldots, J$. It is well known, cp. [25, Equations (3.7) and (3.8)], that for every $n \in \mathbb{N}$ the error that results from one step of this scheme can be written for $\mathbb{P}$-a.e. ω as

$$u(\omega) - U^n(\omega) = E_J(\omega)(u(\omega) - U^{n-1}(\omega)),$$

where the random residual operator is given by

$$E_J = (\mathrm{Id} - T_J) \cdots (\mathrm{Id} - T_1).$$

For a symmetrized version, we refer to [25, Algorithm 3.4], where the respective residual operator E_J^s satisfies that $E_J^s = E_J^* E_J$, where E_J^* denotes the adjoint of E_J with respect to $(\cdot, \cdot)_A$. We observe that $\|E_J^s\|_A = \|E_J\|_A^2$, which implies that convergence for either of the algorithms implies convergence respectively for the other as well. Note that the symmetrized version that results from E_J^s can also be brought into the from of (5), where for $f \in V$

$$Bf := U^1,$$

where U^1 is defined to be the result of the symmetrized version of Algorithm 2, cp. [25, Algorithm 3.4] with $U^0 = 0$. This random symmetric operator can also be used as a preconditioner and will be denotes by B^s.

A multilevel iteration can be defined under the assumption that there are nested subspaces satisfying

$$M_1 \subset M_2 \subset \cdots \subset M_J = V, \tag{9}$$

where we also define operators $\hat{Q}_j, \hat{P}_j : M_J \to M_j$ and $\hat{A}_j : M_j \to M_j$ for $j = 1, \ldots, J$ respectively. The random multilevel iterations $\hat{B}_j^s : M_j \to M_j$, $j = 1, \ldots, J$ with parameters $m, k \in \mathbb{N}$ will be defined iteratively.

Definition 1 Let $m, k \in \mathbb{N}$ and set $\hat{B}_1^s := \hat{A}_1^{-1}$ and assume that $\hat{B}_{j-1}^s : M_{j-1} \to M_{j-1}$ is already defined, then set $v_0 := 0$, $w_0 := 0$ for $v_0, w_0 \in M_j$ and define for every $g \in M_j$ the multigrid iteration $\hat{B}_j^s g$ by:

1. $v_\ell := v_{\ell-1} + \hat{R}_j(g - \hat{A}_j v_{\ell-1})$, for $\ell = 1, \ldots, m$

2. $v_{m+1} := v_m + w_k$, where for $i = 1, \ldots, k$

$$w_i := w_{i-1} + \hat{B}_{j-1}^s[\hat{Q}_{j-1}(g - \hat{A}_j w_{i-1}) - \hat{A}_{j-1} w_{i-1}]$$

3. $\hat{B}_j^s g := v^{2m+1}$, where v^ℓ is given by step 1 for $\ell = m+2, \ldots, 2m+1$.

Algorithm 3 Let U^0 be given, then U^n is defined by the linear iteration in (5) with $B = \hat{B}_J^s$.

Note that for $k = 1$, we have the V-cycle and for $k = 2$ the W-cycle iteration. We will consider multiple and single smoothing versions, i.e., $m \in \mathbb{N}$ and $m = 1$. According to [6, Equation (2.14)] the residual operator for $m = k = 1$ is given by

$$\hat{E}_J^s = (\mathrm{Id} - \hat{B}_J^s \hat{A}_J) = (\mathrm{Id} - \hat{T}_J) \cdots (\mathrm{Id} - \hat{T}_1)(\mathrm{Id} - \hat{T}_1^*) \cdots (\mathrm{Id} - \hat{T}_J^*),$$

where $\hat{T}_j := \hat{R}_j \hat{A}_j \hat{P}_j$. There is also a non symmetric versions of this algorithm, where either step 1 or step 3 are omitted. For $m = k = 1$, the residual operator for this algorithm in the case of omitting step 1 is according to [25, Equation (3.11)] given by

$$\hat{E}_J = (\mathrm{Id} - \hat{B}_J \hat{A}_J) = (\mathrm{Id} - \hat{T}_J) \cdots (\mathrm{Id} - \hat{T}_1),$$

which implies $\hat{E}_J^s = \hat{E}_J \hat{E}_J^*$. Again, it holds that $\|\hat{E}_J^s\|_A = \|\hat{E}_J\|_A^2$. The respective operators will be denoted by $\hat{B}_j$, $j = 1, \ldots, J$.

Algorithm 4 Let U^0 be given, then U^n is defined by the linear iteration in (5) with $B = \hat{B}_J$.

The case where step 3 is omitted is similar and is for example given in [7, Section 2].

For $m = k = 1$, Algorithms 2 and 4 are related, cp. [25] and will be made precise in the following proposition as a version of [25, Propositions 3.2 and 3.3].

Proposition 2.2 *In the framework of Algorithm 4 resulting from given nested subspaces (9), Algorithm 2 results with $V_j = M_j$ and $R_j = \hat{R}_j$, $j = 1, \ldots, J$.*

In the framework of Algorithm 2 resulting from given nested subspaces (7), Algorithm 4 results with $\hat{M}_j := \sum_{i=1}^j V_i$ and $\hat{R}_j = R_j Q_j$, $j = 1, \ldots, J$.

For the convergence of Algorithm 1 we have to prove bounds of $\kappa(BA)$, whereas for Algorithms 2 and 3 we have to show that $\|E_J\|_A \leq \delta$ for some random variable δ taking values in $(0,1)$ $\mathbb{P}$-a.s. According to Proposition 2.2, for $m = k = 1$ it is sufficient to prove the contraction property for the setting of Algorithm 2. Note that also the convergence of the symmetrized versions of the algorithms will be implied. For a brief overview how the presented algorithms are intented to be applied in the ensuing analysis, the additive preconditioner in Algorithm 1 will in later applications be the BPX or additive Schwarz preconditioner. Algorithm 2 can be used as a multiplicative domain decomposition method and Algorithms 3 and 4 refer to non-symmetric and symmetric multigrid.

In [25], assumptions are introduced involving two parameters K_0 and K_1 that allow to discuss convergence in this abstact setting. Here K_0 and K_1 are positive random variables. We recall the two conditions [25, Equations (4.2) and (4.3)].

Assumption 1 There exist positive random variables K_0 and K_1 such that

1. for every $v \in V$, there exists a decomposition $v = \sum_{j=1}^{J} v_j$ with $v_j \in V_j$, $j = 1, \ldots, J$, such that for $\mathbb{P}$-a.e. ω

$$\sum_{j=1}^{J} (R_j^{-1}(\omega)v_j, v_j) \leq K_0(\omega)(A(\omega)v, v) \tag{10}$$

2. for $\mathbb{P}$-a.e. ω and for every $S \subset \{1, \ldots, J\} \times \{1, \ldots, J\}$ and $v_j, w_j \in V$, $j = 1, \ldots, J$,

$$\sum_{(i,j) \in S} (T_i(\omega)v_i, T_j(\omega)w_j)_{A(\omega)}$$

$$\leq K_1(\omega) \left(\sum_{i=1}^{J} (T_i(\omega)v_i, v_i)_{A(\omega)} \right)^{1/2} \left(\sum_{j=1}^{J} (T_j(\omega)w_j, w_j)_{A(\omega)} \right)^{1/2}. \tag{11}$$

Theorem 2.3 *Let Assumption 1 be satisfied. Let B^a be the random preconditioner given by (8). Then, for $\mathbb{P}$-a.e. ω*

$$\kappa(B^a(\omega)A(\omega)) \leq K_0(\omega)K_1(\omega). \tag{12}$$

The residual operator E_J from Algorithm 2 satisfies for $\mathbb{P}$-a.e. ω

$$\|E_J(\omega)\|_{A(\omega)}^2 \leq 1 - \frac{2 - \nu}{K_0(\omega)(1 + K_1(\omega))^2}, \tag{13}$$

where $\nu \geq \max_{j=1,\ldots,J}\{\rho(R_j(\omega)A_j(\omega))\}$ is deterministic.

Proof. These are explicitly [25, Theorems 4.1 and 4.4]. □

Remark 2.4 *Note that the convergence of Algorithms 3 and 4 also holds with the same random contraction number due to Proposition 2.2. The convergence of Algorithms 3 and 4 in the general cases $m, k \in \mathbb{N}$ also holds with this random contraction number respectively with its square root, cp. [7, Equations (2.12) and (2.13)]. But in the setting of Assumption 1 and Theorem 2.3 multiple smoothing or correction steps seem not to improve the contraction number.*

The random parameters K_0 and K_1 can be estimated in some cases with [25, Lemmas 4.5, 4.6 and 4.7]. Let us state a specific case of [25, Lemma 4.6] as the following lemma.

Lemma 2.5 *Let K_3 be a positive random variable that is independent of J and let $\gamma \in (0,1)$ be deterministic. If for $\mathbb{P}$-a.e. ω and for every $v, w \in V$ and every $i, j \in \{1, \ldots, J\}$ it holds that*

$$(T_i(\omega)v, T_j(\omega)w)_{A(\omega)} \leq K_3(\omega)\nu\gamma^{|i-j|}(T_i(\omega)v, v)_{A(\omega)}^{1/2}(T_j(\omega)w, w)_{A(\omega)}^{1/2}, \tag{14}$$

then

$$K_1(\omega) \leq K_3(\omega)\nu\frac{2}{1-\gamma}.$$

Note that (14) is often called *strengthened Cauchy–Schwarz inequality*.

The other approach to prove convergence of multigrid methods as Algorithms 3 and 4 takes advantage of multiple smoothing steps, i.e., $m \in \mathbb{N}$, and relies on the so called *regularity and approximation* assumption, which seems to have been introduced in this form by Bramble and Pasciak in [4].

Assumption 2 Let $s \in (0,1]$. There exist positive random variables C_R and C_s such that:

1. for every $j = 1, \ldots, J$ and for $\mathbb{P}$-a.e. ω

$$\frac{\|v\|^2}{\lambda_{\max}(\hat{A}_j(\omega))} \leq C_R(\omega)(\hat{R}_j(\omega)v, v) \quad \forall v \in V_j.$$

2. for every $j = 1, \ldots, J$ and for $\mathbb{P}$-a.e. ω

$$((\mathrm{Id} - \hat{P}_{j-1}(\omega))v, v)_{A(\omega)} \leq C_s(\omega)\left(\frac{\|\hat{A}_j(\omega)v\|^2}{\lambda_{\max}(\hat{A}_j(\omega))}\right)^s (v, v)_{A(\omega)}^{1-s} \quad \forall v \in M_j.$$

The convergence estimate for Algorithm 3 for general $m \in \mathbb{N}$ and $s \in (0,1)$ was proven in [4], whereas the result for $s = 1$ has been known earlier, cp. [17]. The following theorem is a version of [4, Theorem 3, Theorem 1, Equation (3.28)]. There increasing the number of smoothing steps m improves the contraction number.

Theorem 2.6 *Let Assumption 2 be satisfied. For $k = 2$, $m \in \mathbb{N}$, and the random variable $C_W := 2^{(1-s)/s}s(1-s)^{(1-s)/s}C_s^{2/s}C_R$ it holds for $\mathbb{P}$-a.e. ω*

$$\|\mathrm{Id} - \hat{B}_J^s(\omega)A(\omega)\|_{A(\omega)} \leq \frac{1}{(1 + m/C_W(\omega))^s}.$$

For $s = 1$, $k = 1$, $m \in \mathbb{N}$, and the random variable $C_V := C_s^2 C_R/2$ it holds for $\mathbb{P}$-a.e. ω

$$\|\mathrm{Id} - \hat{B}_J^s(\omega)A(\omega)\|_{A(\omega)} \leq \frac{1}{1 + m/C_V(\omega)}.$$

The improvement of the contraction number with increased smoothing steps may be interesting in some application. For convergence estimates for Algorithm 4 for general $m \in \mathbb{N}$ we refer to [4, Theorem 4]. There the contraction number is the square root of the one given in the previous theorem. For the case of one correction step, i.e., $k = 1$, only non uniform convergence estimates seem to be available under Assumption 2, cp. [4, Theorems 1 and 2].

Proposition 2.7 *Let $B \in \{B^s, \hat{B}_j^s\}$ and let δ be a random variable taking values in $(0,1)$. If for $\mathbb{P}$-a.e. ω, $\|\mathrm{Id} - B(\omega)A(\omega)\|_{A(\omega)} \leq \delta(\omega)$, then $\kappa(B(\omega)A(\omega)) \leq 1/(1-\delta(\omega))$.*

Proof. Since the operator $\mathrm{Id} - BA$ can be written as E^*E for some appropriate E that was discussed above, it holds that $((\mathrm{Id} - BA)v, v)_A = \|E^*v\|_A^2 \geq 0$ for every $v \in V$. Hence,

$$(BAv, v)_A \leq (v, v)_A \quad \forall v \in V. \tag{15}$$

The assumption implies that

$$(1 - \delta)(v, v)_A \leq (BAv, v)_A \quad \forall v \in V, \tag{16}$$

which then implies the claim. $\qquad\square$

Remark 2.8 *The operators B^s and $\hat{B}_J^s$ can also be used as preconditioners in a PCG method to accelerate convergence. The respective condition number can be bounded by Proposition 2.7 in combination with Theorem 2.3 or Theorem 2.6.*

3   Strong convergence of iterative methods

We recall the possibly infinite dimensional Hilbert space $\mathcal{V}$ and let $(\mathcal{H}, (\cdot, \cdot))$ be another Hilbert spaces such that the embedding $\mathcal{H} \subset \mathcal{V}$ is continuous. Let $(\mathcal{V}_\ell, \ell \in \mathbb{N})$ be a nested sequence of finite dimensional subspaces of $\mathcal{V}$, i.e., $\mathcal{V}_1 \subset \mathcal{V}_2 \subset \cdots \subset \mathcal{V}$. Let the finite dimensional spaces $\mathcal{V}_\ell$ have dimensions $N_\ell := \dim(\mathcal{V}_\ell)$, $\ell \in \mathbb{N}$. Similar to Section 2, we introduce random operators $\mathcal{A}_\ell$, $\mathcal{P}_\ell$, R_ℓ, T_ℓ, and Q_ℓ with respect to the inner product $(\cdot, \cdot)$ of $\mathcal{H}$, $\ell \in \mathbb{N}$. The "usual inner product" $(\cdot, \cdot)_{\mathbb{E}(\mathcal{A})}$ on $\mathcal{V}$ is given by

$$(v, w)_{\mathbb{E}(\mathcal{A})} := \langle \mathbb{E}(\mathcal{A})v, w \rangle_{\mathcal{V}^*, \mathcal{V}} \quad \forall v, w \in \mathcal{V}.$$

The inner product $(\cdot, \cdot)_{\mathcal{A}}$ with respect to the random symmetric operator $\mathcal{A}$ will also be considered.

For every $\ell \in \mathbb{N}$, we consider the variational form of (1) on the suchspace $\mathcal{V}_\ell$. For every $\ell \in \mathbb{N}$, this gives rise to the random linear equation: find $u_\ell : \Omega \to \mathcal{V}_\ell$ such that for $\mathbb{P}$-a.e. ω

$$_{\mathcal{V}^*}\langle \mathcal{A}_\ell(\omega)u_\ell(\omega), v_\ell \rangle_{\mathcal{V}} = {}_{\mathcal{V}^*}\langle f, v_\ell \rangle_{\mathcal{V}}, \quad \forall v_\ell \in \mathcal{V}_\ell,$$

This is uniquely solvable by the Lax–Milgram lemma.

We assume that the random solution $u : \Omega \to \mathcal{V}$ of the problem in (1) is approximated by the Galerkin approximations u_ℓ, $\ell \in \mathbb{N}$. Specifically, we assume that there exists a positive random variable $\mathcal{C}$ that may depend on u, a scalar $s > 0$, and a decreasing sequence $(h_\ell : \ell \in \mathbb{N})$ such that for every $\ell \in \mathbb{N}$ and for $\mathbb{P}$-a.e. ω

$$\|u(\omega) - u_\ell(\omega)\|_{\mathcal{A}(\omega)} \leq \mathcal{C}(\omega)h_\ell^s. \tag{17}$$

We apply an iterative method such as Algorithm 1, 2, 3, or 4 with random contraction number δ, that is independent of ℓ, with n iterations and solve exactly on level 1, i.e., starting with $U_{\ell-1}^n$ as initial guess for level ℓ we carry out n iterations of the algorithm with a random contraction number δ that takes values in $(0, 1)$. Hence, we obtain a sequence U_ℓ^n, $\ell \in \mathbb{N}$, where we set $U_1^n := u_1$. This multilevel process was used in [1, Section 3] to derive optimal complexity bounds for the solution of certain linear equations and is also commonly referred to as "full multigrid".

Lemma 3.1 *Let us assume that $(u_\ell, \ell \in \mathbb{N})$ satisfies (17) for some $s > 0$ and let the sequence $(h_\ell, \ell \in \mathbb{N})$ satisfy that $h_\ell = \gamma^\ell h_0$, $\ell \in \mathbb{N}$, for some fixed $h_0 > 0$ and $\gamma \in (0, 1)$. Then, for every $\ell \in \mathbb{N}$, U_ℓ^n and for $\mathbb{P}$-a.e. ω*

$$\|u(\omega) - U_\ell^n(\omega)\|_{\mathcal{A}(\omega)} \leq 2\mathcal{C}(\omega)\left(\sum_{k=0}^{\ell-1} h_{\ell-k}^s \delta(\omega)^{nk}\right) = 2\mathcal{C}(\omega)\left(\sum_{k=0}^{\ell-1}(\gamma^{-s}\delta(\omega)^n)^k\right)h_\ell^s.$$

Proof. The argument is for example given in [3, Chapter 10]. We recall that on every level $\ell \geq 2$ we will carry out n iterations of our contractive algorithm with initial guess $U_{\ell-1}^n$. Thus,

$$\|u_\ell - U_\ell^n\|_\mathcal{A} \leq \delta^n \|u_\ell - U_{\ell-1}^n\|_\mathcal{A} \quad \ell = 2, \ldots, J.$$

A twofold application of the triangle inequality implies with (17) a recursion, i.e.,

$$\begin{aligned}
\|u - U_\ell^n\|_\mathcal{A} &\leq \|u - u_\ell\|_\mathcal{A} + \|u_\ell - U_\ell^n\|_\mathcal{A} \\
&\leq \|u - u_\ell\|_\mathcal{A} + \delta^n \|u_\ell - U_{\ell-1}^n\|_\mathcal{A} \\
&\leq \|u - u_\ell\|_\mathcal{A}(1 + \delta^n) + \delta^n(\|u - U_{\ell-1}^n\|_\mathcal{A}) \\
&\leq \mathcal{C}h_\ell^s(1 + \delta^n) + \delta^n\|u - U_{\ell-1}^n\|_\mathcal{A}.
\end{aligned}$$

Since we have set $U_1^n := u_1$, this resursion terminates and can be expanded as claimed using that $\max\{\mathcal{C}(1 + \delta^n), \mathcal{C}\} \leq 2\mathcal{C}$, where the second claim follows with the assumption that $h_\ell = \gamma^\ell h_0$, $\ell \in \mathbb{N}$. $\qquad\square$

Theorem 3.2 *Let the assumptions of Lemma 3.1 be satisfied. Let us assume that $(\mathcal{C}/\sqrt{\check{a}}) \in L^p(\Omega)$ for some $p \in [1, +\infty)$ and that $1/(1-\delta) \in L^{p'}(\Omega)$ for some $p' \in [1, +\infty)$. For every $\eta \in (0, 1)$, every deterministic number of iterations $n \in \mathbb{N}$, $q \in [1, p]$, and $r := p'(p-q)/(pq)$ it holds that for every $\ell \geq 2$*

$$\|u - U_\ell^n\|_{L^q(\Omega;\mathcal{V})} \leq \mathcal{C}_{q,p,p',s,\eta} \left(h_\ell^s + n^{-r}\right),$$

where

$$\mathcal{C}_{q,p,p',s,\eta} := 2\max\left\{\frac{\|\mathcal{C}/\sqrt{\check{a}}\|_{L^q(\Omega)}}{1-\eta}, \|\mathcal{C}/\sqrt{\check{a}}\|_{L^p(\Omega)}\frac{h_1^s}{1-\gamma^s}\left\|\frac{1}{1-\delta}\right\|_{L^{p'}(\Omega)}^r \log\left(\frac{1}{\eta\gamma^s}\right)^r\right\}.$$

If $p' = +\infty$ and $p \in [1, +\infty]$, then for every $n \in \mathbb{N}$ satisfying $\|\delta\|_{L^\infty(\Omega)}^n \leq \eta\gamma^s$ and every $q \in [1, p]$, it holds that for every $\ell \geq 2$

$$\|u - U_j^n\|_{L^q(\Omega;\mathcal{V})} \leq 2\|\mathcal{C}/\sqrt{\check{a}}\|_{L^q(\Omega)}\frac{1}{1-\eta}h_\ell^s.$$

Proof. The idea of the proof is to decompose the probability space into $\Omega = \Omega_n \cup (\Omega_n)^\complement$, where

$$\Omega_n := \{\omega \in \Omega : \delta(\omega)^n < \eta\gamma^s\}$$

and therefore $(\Omega_n)^\complement = \{\omega \in \Omega : \delta(\omega)^n \geq \eta\gamma^s\}$. Note that both sets are measurable. For notational convenience, we omit ω in the following when discussing subsets of Ω. Our goal is to show, how the probability of $(\Omega_n)^\complement$ tends to zero for increasing values of $n \in \mathbb{N}$, to be able to justify the applicability of a well known argument on the sets Ω_n, $n \in \mathbb{N}$. Naturally, $\Omega_{n_1} \subset \Omega_{n_2}$ for every choice of natural numbers $n_1 \leq n_2$. We recall a version of the Markov inequality, cp. [2, Equation (2.1)], i.e., for a random variable X taking values in $(1, +\infty)$ and a non-decreasing function ϕ such that $\phi(t) > 0$ it holds that

$$\mathbb{P}(X \geq t) \leq \mathbb{P}(\phi(X) \geq \phi(t)) \leq \frac{\mathbb{E}(\phi(X))}{\phi(t)}, \quad t \in (1, +\infty).$$

We select the non decreasing positive function $\phi(t) := t^{p'}$. Then, for $X = 1/(1-\delta)$ and every $t \in (1, +\infty)$ it holds that

$$\mathbb{P}\left(\frac{1}{1-\delta} \geq t\right) \leq \mathbb{E}\left(\left(\frac{1}{1-\delta}\right)^{p'}\right)\frac{1}{t^{p'}}. \tag{18}$$

Since for every $r \in (0,1)$ the equality of measurable sets $\{\delta \geq r\} = \{1/(1-\delta) \geq 1/(1-r)\}$ holds, we conclude that

$$\mathbb{P}((\Omega_n)^{\mathsf{c}}) = \mathbb{P}\left(\delta^n \geq \eta\gamma^s\right) = \mathbb{P}\left(\delta \geq (\eta\gamma^s)^{1/n}\right) = \mathbb{P}\left(\frac{1}{1-\delta} \geq \frac{1}{1-(\eta\gamma^s)^{1/n}}\right).$$

We observe that the function $x \mapsto (1 - (\eta\gamma^s)^{1/x})x$ from $(1, +\infty)$ to $(1, +\infty)$ is increasing. Since the well known rule of L'Hospital implies that $\lim_{x \to +\infty}(1 - (\eta\gamma^s)^{1/x})x = \log(1/(\eta\gamma^s))$, we conclude that for every $n \in \mathbb{N}$

$$1 - (\eta\gamma^s)^{1/n} \leq \log\left(\frac{1}{\eta\gamma^s}\right)\frac{1}{n}.$$

For every $n \in \mathbb{N}$, we choose $t := 1/(1 - (\eta\gamma^s)^{1/n})$ in (18), and conclude that for every $n \in \mathbb{N}$ it holds that

$$\mathbb{P}((\Omega_n)^{\mathsf{c}}) \leq \mathbb{E}\left(\left(\frac{1}{1-\delta}\right)^{p'}\right)\log\left(\frac{1}{\eta\gamma^s}\right)^{p'}\left(\frac{1}{n}\right)^{p'}. \tag{19}$$

Hence, we have established estimates for the probability of the sets $(\Omega_n)^{\mathsf{c}}$, $n \in \mathbb{N}$. We apply Lemma 3.1, (2), and decompose the probability space into $\Omega = \Omega_n \cup (\Omega_n)^{\mathsf{c}}$ to obtain that for q as in the statement of the theorem

$$\mathbb{E}(\|u - U_\ell^n\|_{\mathbb{E}(\mathcal{A})}^q) \leq \mathbb{E}\left(\left(2\frac{\mathcal{C}}{\sqrt{\tilde{a}}}\sum_{k=0}^{\ell-1}h_{\ell-k}^s\delta^{nk}\right)^q\right)$$

$$= \mathbb{E}\left(\left(2\frac{\mathcal{C}}{\sqrt{\tilde{a}}}\sum_{k=0}^{\ell-1}(\gamma^{-s}\delta^n)^k\right)^q\mathbb{1}_{\Omega_n}\right)h_\ell^{sp}$$

$$+ \mathbb{E}\left(\left(2\frac{\mathcal{C}}{\sqrt{\tilde{a}}}\sum_{k=0}^{\ell-1}h_{\ell-k}^s\delta^{nk}\right)^q\mathbb{1}_{(\Omega_n)^{\mathsf{c}}}\right).$$

Since on Ω_n holds that $\gamma^{-s}\delta^n < \eta$, we obtain with a geometric series argument that

$$\mathbb{E}\left(\left(2\frac{\mathcal{C}}{\sqrt{\tilde{a}}}\sum_{k=0}^{\ell-1}(\gamma^{-s}\delta^n)^k\right)^q\mathbb{1}_{\Omega_n}\right)h_\ell^{sq} \leq 2^q\mathbb{E}\left(\left(\frac{\mathcal{C}}{\sqrt{\tilde{a}}}\right)^q\right)\left(\frac{1}{1-\eta}\right)^q h_\ell^{sq}.$$

The relation $h_\ell = \gamma^\ell h_0$, $\ell \in \mathbb{N}$, implies that for every $\ell \geq 2$ it holds that $\sum_{k=0}^{\ell-1}h_{\ell-k}^s \leq \sum_{\ell \geq 1}\gamma^{\ell s}h_0 = h_1^s/(1-\gamma^s)$. The Hölder inequality with $r_1 = p/q$ and $r_2 = p/(p-q)$ implies with the tail bound of δ^n in (19) that

$$\mathbb{E}\left(\left(2\frac{\mathcal{C}}{\sqrt{\tilde{a}}}\sum_{k=0}^{\ell-1}h_{\ell-k}^s\delta^{nk}\right)^q\mathbb{1}_{(\Omega_n)^{\mathsf{c}}}\right)$$

$$\leq 2^q\left\|\frac{\mathcal{C}}{\sqrt{\tilde{a}}}\right\|_{L^p(\Omega)}^q\left(\frac{h_1^s}{1-\gamma^s}\right)^q\mathbb{E}\left(\mathbb{1}_{(\Omega_n)^{\mathsf{c}}}\right)^{1/r_2}$$

$$\leq 2^q\left\|\frac{\mathcal{C}}{\sqrt{\tilde{a}}}\right\|_{L^p(\Omega)}^q\left(\frac{h_1^s}{1-\gamma^s}\right)^q\left\|\frac{1}{1-\delta}\right\|_{L^{p'}(\Omega)}^{rq}\log\left(\frac{1}{\eta\gamma^s}\right)^{rq}\left(\frac{1}{n}\right)^{rq},$$

which finishes the proof of the theorem for $p' \in [1, +\infty)$.

If $p' = +\infty$ and therefore $1/(1-\delta) \in L^\infty(\Omega)$, then $\mathbb{P}$-a.s. it holds that

$$\delta \leq \left(\left\|\frac{1}{1-\delta}\right\|_{L^\infty(\Omega)} - 1\right)\left\|\frac{1}{1-\delta}\right\|_{L^\infty(\Omega)}^{-1} < 1.$$

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Consequently, $\|\delta\|_{L^\infty(\Omega)} < 1$ and there exists a deterministic $n \in \mathbb{N}$ such that $\mathbb{P}$-a.s. $\gamma^{-s}\delta^n \leq \gamma^{-s}\|\delta\|_{L^\infty(\Omega)}^n < \eta$. Hence, the geometric series argument that was carried out on the set Ω_n can be applied verbatim on Ω, which implies the claim. $\qquad\square$

Remark 3.3 *Let X be a random variable with values in $(0,1)$ such that $1/(1-X) \in L^p(\Omega)$ for some $p \in [1, +\infty)$, then (19) implies that for every $\eta < 1$ and $n \in \mathbb{N}$*

$$\mathbb{P}(X^n \geq \eta) \leq \left\|\frac{1}{1-X}\right\|_{L^p(\Omega)}^p \log\left(\eta^{-1}\right)^p n^{-p}.$$

Note that the argument to bound the error contributions from the different levels of the multilevel method independently from the number of levels in the deterministic case can for example be found in [3, Chapter 10, p. 130].

The convergence estimate (17) may not be available. Nevertheless, the approximated solution by an iterative scheme still converges strongly to the discrete solution on fixed levels. Let U_ℓ^n be the result of n steps of an iterative scheme with random contraction number δ_ℓ taking values in $(0,1)$ to approximate u_ℓ with initial guess $0 \in \mathcal{V}_\ell$, $\ell \in \mathbb{N}$, i.e., for $\mathbb{P}$-a.e. ω

$$\|u_\ell(\omega) - U_\ell^n(\omega)\|_{\mathcal{A}(\omega)} \leq \delta_\ell(\omega)^n \|u_\ell(\omega)\|_{\mathcal{A}(\omega)}.$$

Note that we allow here the random contraction number to depend on ℓ.

Theorem 3.4 *Let us assume that $[\omega \mapsto \|u_\ell(\omega)\|_{\mathbb{E}(\mathcal{A})}/\sqrt{\check{a}(\omega)}] \in L^p(\Omega; V)$ for some $p \in [1, +\infty)$ and that $1/(1-\delta_\ell) \in L^{p'}(\Omega)$ for some $p' \in [1, +\infty)$. For every $\eta \in (0,1)$, every deterministic number of iterations $n \in \mathbb{N}$, $q \in [1, p]$, and $r := p'(p-q)/(pq)$ it holds that for every $\ell \geq 1$*

$$\|u_\ell - U_\ell^n\|_{L^q(\Omega; \mathcal{V})} \leq \mathcal{C}_{q,p,p',s,\eta}(\eta^n + n^{-r}).$$

where

$$\mathcal{C}_{q,p,p',s,\eta} := 2 \max\left\{ \left\|\frac{\|u_\ell\|_{\mathbb{E}(\mathcal{A})}}{\sqrt{\check{a}}}\right\|_{L^p(\Omega)}, \left\|\frac{1}{1-\delta_\ell}\right\|_{L^{p'}(\Omega)}^r \log\left(\eta^{-1}\right)^r \right\}.$$

If $p' = +\infty$ and $p \in [1, +\infty]$, then for every $n \in \mathbb{N}$ satisfying $\|\delta\|_{L^\infty(\Omega)}^n \leq \eta\gamma^s$ and every $q \in [1, p]$, it holds that for every $\ell \in \mathbb{N}$

$$\|u_\ell - U_\ell^n\|_{L^q(\Omega; \mathcal{V})} \leq \left\|\frac{\hat{a}}{\check{a}}\|u_\ell\|_V\right\|_{L^p(\Omega)} \|\delta_\ell\|_{L^\infty(\Omega)}^n.$$

Proof. The proof is similar to the proof of Theorem 3.2 using Remark 3.3. $\qquad\square$

4 Multigrid methods

Here, we provide sufficient conditions under which the regularity and approximation assumption and the strengthened Cauchy–Schwarz inequality can be proven with explicit dependence on the operator $\mathcal{A}$. This will allow us to show $L^q(\Omega)$ bounds of the condition numbers and tail bounds of the random contraction number in order to apply the strong error bounds from Theorems 3.2 and 3.4.

4.1 Verification of Assumption 1

We turn to a proof of the strengthened Cauchy–Schwarz inequality. We will provide a proof in the case of a random, symmetric elliptic differential operators. To be specific, let $H_0^1(D)$, $H^s(D)$, $s \in [-1, 2]$, be the usual Sobolev spaces for some polytopal domain $D \subset \mathbb{R}^d$, $d \geq 1$ arbitrary, such that $H^{-s}(D)$ is the dual space of $H^s(D)$, $s \in [0, 1]$, and $H^0(D) = L^2(D)$. The reader is referred to [16, Chapter 1] for details on Sobolev spaces. We consider the class of random symmetric operators

$$\mathcal{A} := -\sum_{i,j=1}^{d} \frac{\partial}{\partial x_i} \left(a_{ij} \frac{\partial}{\partial x_j} \right) + a : H_0^1(D) \to (H_0^1(D))^*,$$

where $(a_{ij}(x))_{i,j=1,\dots,d}$ is a random symmetric matrices for a.e. $x \in D$. Let us assume that $\overline{D} = \bigcup_{k=1}^{K} \overline{D}_k$, where the subdomains D_k are pairwise disjoint with a polytopal boundary. Furthermore, we assume that the random fields a and a_{ij} are strongly measurable as mappings from Ω to $L^\infty(D)$, and $a_{ij}|_{D_k}$ is strongly measurable as a mapping from Ω to $W^{s,p}(D_k)$ such that $s > d/p$, $k = 1, \dots, K$, $i, j = 1, \dots, d$, where $W^{s,p}(D_k)$, $s \geq 0$, $p \in [1, +\infty)$, $k = 1, \dots, K$, denote the Sobolev–Slobodeckij spaces, cp. [16, Definition 1.3.2.1]. We assume that for $\mathbb{P}$-a.e. ω, $\operatorname{ess\,inf}_{x\in D}\{a(\omega, x)\} \geq 0$ and that there exists a strictly positive random variable $\check{a}$ such that for $\mathbb{P}$-a.e. ω

$$\operatorname*{ess\,inf}_{x\in D} \left\{ \sum_{i,j=1}^{d} a_{ij}(\omega, x)\xi_i\xi_j \right\} \geq \check{a}(\omega)|\xi|^2 \quad \forall \xi \in \mathbb{R}^d.$$

The corresponding bilinear form $a(\cdot, \cdot)$ is given by

$$a(v, w) = \sum_{i,j=1}^{d} \int_D a_{ij} \frac{\partial v}{\partial x_j} \frac{\partial w}{\partial x_i} + avw\mathrm{d}x \quad \forall v, w \in H_0^1(D).$$

Furthermore, let $\hat{a}, \check{a}$ be strictly positive random variables such that for $\mathbb{P}$-a.e. ω

$$\check{a}(\omega) \int_D |\nabla v|^2 \mathrm{d}x \leq a_\omega(v, v) \leq \hat{a}(\omega) \int_D |\nabla v|^2 \mathrm{d}x \quad \forall v \in H_0^1(D). \tag{20}$$

The following example will is class of random coefficients, which will be further discussed in Section 6.

Example 1 The class of lognormal random coefficient fields $a_{ij} = \exp(Z)\delta_{ij}$ and $a = 0$, where $Z : \Omega \to W^{s,p}(D)$ is a strongly measurable Gaussian random field and δ_{ij} denotes the Kronecker symbol, satisfy these conditions.

The proof of the strengthened Cauchy–Schwarz inequality draws on [6, Sections 4 and 5] and [25, Section 4]. The setting in reference [6] allows for low Hölder regularity of the coefficients of elliptic operators, but does not provide a strengthened Cauchy–Schwarz inequality needed for our setting of Assumption 1. The strengthened Cauchy–Schwarz inequality proved in [25, Lemmas 6.1 and 6.3] is limited to coefficients with $W^{1,\infty}(D)$ regularity. Here, a strengthened Cauchy–Schwarz inequality will be proved with explicit dependence on the coefficients that is valid for arbitrary low Hölder regularity of the coefficients and also allows for jumps across ∂D_k (see ahead Proposition 4.2). We will identify estimates with explicit dependence on the random coefficients of $\mathcal{A}$.

Let $(\mathcal{T}_\ell, \ell \in \mathbb{N})$ be a nested sequence of shape regular simplicial, uniformly refined meshes of D, i.e., every $\tau \in \mathcal{T}_\ell$ is a finite union of elements in $\mathcal{T}_{\ell+1}$, $\ell \in \mathbb{N}$. Note that in one refinement step one simplex is refined into 2^d subsimplices. For every $k = 1, \dots, K$ and $\ell \in \mathbb{N}$, we require that

$\overline{D}_k = \bigcup_{\tau \in \mathcal{T}_\ell, D_k \cap \tau \neq \emptyset} \overline{\tau}$. Let $\mathcal{V}_\ell \subset \mathcal{V}$ be the space of piecewise polynomial function with respect to the mesh $\mathcal{T}_\ell$, $\ell \in \mathbb{N}$. For simplicity, we will consider here only first order FE, i.e., polynomial degree one. Define $-\Delta_\ell : \mathcal{V}_\ell \to \mathcal{V}_\ell$ by the bilinear form $(w_\ell, v_\ell) \mapsto \int_D \nabla w_\ell \cdot \nabla v_\ell \mathrm{d}x$ over $\mathcal{V}_\ell \times \mathcal{V}_\ell$, $\ell \in \mathbb{N}$. By [6, Equation (5.1)], there exists a deterministic C such that for every $j \in \mathbb{N}$ and every $v \in \mathcal{V}_j$ there exists a decomposition $v = \sum_{i=1}^{j} v_i$ with $v_i \in \mathcal{V}_i$ such that

$$\sum_{i=1}^{j} \lambda_{\max}(-\Delta_i) \|v_i\|_{L^2(D)}^2 \leq C \int_D |\nabla v|^2 \mathrm{d}x. \tag{21}$$

Moreover, the following inverse estimates hold: there exists a constant C such that for every $v \in \mathcal{V}_\ell$

$$\|v\|_{H^1(D)} \leq C h_\ell^{-1} \|v\|_{L^2(D)} \quad \text{and} \quad \|v\|_{H^{1+s}(D)} \leq C h_\ell^{-s} \|v\|_{H^1(D)}, \tag{22}$$

where

$$h_\ell := \max_{\tau \in \mathcal{T}_\ell} \{\mathrm{diam}(\tau)\}$$

and $s \in (0, 1/2)$, cp. [13, Theorem 3.2.6] and [9, Equation (10.1)]. These inverse estimate are sharp, which can be seen by choosing v to be a nodal basis function of the FE space $\mathcal{V}_\ell$. Since by (20) for $\mathbb{P}$-a.e. ω

$$\check{a}(\omega)\lambda_{\max}(-\Delta_\ell) \leq \lambda_{\max}(\mathcal{A}_\ell(\omega)) \leq \hat{a}(\omega)\lambda_{\max}(-\Delta_\ell), \tag{23}$$

we also observe that for $\mathbb{P}$-a.e. ω

$$\lambda_{\max}(\mathcal{A}_\ell(\omega)) \geq C\check{a}(\omega)h_\ell^{-2}, \tag{24}$$

which for $-\Delta_\ell$ is a consequence of the sharpness of (22).

We require the following assumptions on the smoothers: there exists a deterministic $\nu \in (0, 2)$ such that for every $j = 1, \ldots, J$, and $\mathbb{P}$-a.e. ω

$$(T_j(\omega)v, T_j(\omega)v)_{\mathcal{A}(\omega)} \leq \nu(T_j(\omega)v, v)_{\mathcal{A}(\omega)} \quad \forall v \in \mathcal{V}_j. \tag{A}$$

There exists deterministic $c_0, c_1 > 0$ such that for every $j = 1, \ldots, J$, and for $\mathbb{P}$-a.e. ω

$$c_0 \frac{\|v\|_{L^2(D)}^2}{\lambda_{\max}(\mathcal{A}_j(\omega))} \leq (R_j(\omega)v, v) \leq c_1 \frac{\|v\|_{L^2(D)}^2}{\lambda_{\max}(\mathcal{A}_j(\omega))} \quad \forall v \in \mathcal{V}_j. \tag{B}$$

Assume that there exists $\gamma \in (0, 1)$ such that for all $i, j \in \mathbb{N}$ satisfying $i \leq j$ it holds that

$$\frac{h_j}{h_i} \leq \gamma^{j-i}, \tag{C}$$

Note that (A) implies that $\rho(R_i A_i) \leq \nu$ for every i. There exist smoothers that satisfy these assumptions, cp. [10, Chapter 8] and [5, 8].

Lemma 4.1 *Let $s \in (0, 1/2)$ and $p \in (d/s, +\infty)$, then for $\mathbb{P}$-a.e. ω, for every $\eta > 0$, $\phi \in H^1(D)$, and $\psi \in H^{1+s}(D)$ it holds that*

$$|a(\omega)(\phi, \psi)| \leq \mathcal{C}(\omega)(\eta^{-1}\|\phi\|_{L^2(D)}^2 + \eta^{s/(1-s)}\|\phi\|_{H^1(D)}^2)^{1/2}\|\psi\|_{H^{1+s}(D)},$$

where for a deterministic constant C independent of $(a_{ij}, a, i, j = 1, \ldots, d)$

$$\mathcal{C}(\omega) := C\left(\max_{k=1,\ldots,K} \sum_{i,j=1}^{d} \|a_{ij}(\omega)\|_{W^{s,p}(D_k)} + \|a(\omega)\|_{L^\infty(D)}\right).$$

Proof. The following argument originates from the proof of [6, Lemma 4.3]. We will track the dependence on the random elliptic coefficients $(a_{ij}, a, i, j = 1, \ldots, d)$. Let us fix $k \in \{1, \ldots, K\}$. There exists a bounded linear extension operator $I_k : H^1(D_k) \to H^1(\mathbb{R}^d)$, e.g. cp. [16, Theorem 1.4.3.1]. For every function $v : D_k \to \mathbb{R}$ the zero extension to $\mathbb{R}^d$ is denoted by $\widetilde{v}$. Let $\mathscr{F}$ denote the Fourier transform on $\mathbb{R}^d$. We obtain with Plancherel's theorem

$$\int_{D_k} a_{ij} \frac{\partial \phi}{\partial x_i} \frac{\partial \psi}{\partial x_j} \mathrm{d}x = \int_{\mathbb{R}^d} \frac{\partial (I_k \phi)}{\partial x_i} \widetilde{a_{ij} \frac{\partial \psi}{\partial x_j}} \mathrm{d}x = \left(\mathscr{F}\left(\frac{\partial (I_k \phi)}{\partial x_i} \right), \mathscr{F}\left(\widetilde{a_{ij} \frac{\partial \psi}{\partial x_j}} \right) \right)_{L^2(\mathbb{R}^d)}.$$

Recall that $\|(1 + |\xi|^2)^{s/2} \mathscr{F}(v)\|_{L^2(\mathbb{R}^d)}$ is the Bessel potential norm in the Hilbert case of order s of a function v. The fact that in the Hilbert case Bessel potential and Slobodeckij spaces are equal with equivalent norms, cp. [24, Definition 2.3.1(d), Theorem 2.3.2(d), Equation 4.4.1(8)], and the boundedness of the zero extension as an operator from $H^s(D_k)$ to $H^s(\mathbb{R}^d)$, cp. [16, Corollary 1.4.4.5], imply with the Cauchy–Schwarz inequality and differentiation rules for $\mathscr{F}$ that exists a constant C such that

$$\int_{D_k} a_{ij} \frac{\partial \phi}{\partial x_i} \frac{\partial \psi}{\partial x_j} \mathrm{d}x \leq C \left(\int_{\mathbb{R}^d} \frac{|\xi|^2}{(1 + |\xi|^2)^s} \mathscr{F}(I_k \phi) \mathrm{d}\xi \right)^{1/2} \left\| a_{ij} \frac{\partial \psi}{\partial x_j} \right\|_{H^s(D_k)}$$

$$\leq C(\eta^{-1} \|\phi\|_{L^2(D_k)}^2 + \eta^{s/(1-s)} \|\phi\|_{H^1(D_k)}^2)^{1/2} \left\| a_{ij} \frac{\partial \psi}{\partial x_j} \right\|_{H^s(D_k)},$$

where the inequality $|\xi|^2/(1 - |\xi|^2)^s \leq \eta^{-1} + \eta^{s/(1-s)}(1 + |\xi|^2)$ is derived with elementary manipulations for every $\eta > 0$ and every $\xi \in \mathbb{R}^d$. By [16, Theorem 1.4.4.2], the multiplication of elements of $W^{s,p}(D_k)$ is a bounded linear operator on $H^s(D_k)$. Thus, by summing over k and the Cauchy–Schwarz inequality there exists a constant C such that

$$\int_D a_{ij} \frac{\partial \phi}{\partial x_i} \frac{\partial \psi}{\partial x_j} \mathrm{d}x$$

$$\leq C \max_{k=1,\ldots,K} \|a_{ij}\|_{W^{s,p}(D_k)} (\eta^{-1} \|\phi\|_{L^2(D)}^2 + \eta^{s/(1-s)} \|\phi\|_{H^1(D)}^2)^{1/2} \|\psi\|_{H^{1+s}(D)}.$$

Since it also holds that

$$\int_D a\phi\psi \mathrm{d}x \leq \|a\|_{L^\infty(D)} (\eta^{-1} \|\phi\|_{L^2(D)}^2 + \eta^{s/(1-s)} \|\phi\|_{H^1(D)}^2)^{1/2} \|\psi\|_{H^{1+s}(D)},$$

the claim of the lemma follows. $\qquad \square$

Proposition 4.2 *Let Assumptions* (**A**) *and* (**B**) *be satisfied by the smoothers* $(R_j : j = 1 \ldots, J)$, *let Assumption* (**C**) *hold, and let* $s \in (0, 1/2)$ *and* $p \in (d/s, +\infty)$. *Then for some deterministic constant* C *independent of* J *the inequality* (11) *from Assumption 1 holds with the random variable*

$$K_1 := C \left(\max_{k=1,\ldots,K} \sum_{i,j=1}^{d} \|a_{ij}\|_{W_p^s(D_k)} + \|a\|_{L^\infty(D)} \right)^2 \left(\frac{1}{\check{a}} \right)^2.$$

Proof. The proof of this proposition merges ideas of the proofs of [6, Lemma 4.2] and [25, Lemma 6.3] to obtain a strengthened Cauchy–Schwarz inequality with explicit dependence on the coefficients $(a_{ij}, a : i, j = 1, \ldots, d)$ in the setting considered here that allows for low spatial regularity of $(a_{i,j})$, $i, j = 1, \ldots, d$. We may assume that $j \geq i$ due to the symmetry of $(\cdot, \cdot)_{\mathcal{A}}$ and let $w \in \mathcal{V}_i$ and $\phi \in \mathcal{V}_j$ be arbitrary, which are both elements of H^{1+s} due to $s < 1/2$. The first inverse estimate in (22) and Lemma 4.1 imply that

$$|a(w, \phi)| \leq \mathcal{C}(\eta^{-1} + \eta^{s/(1-s)} h_j^{-2})^{1/2} \|\phi\|_{L^2(D)} \|w\|_{H^{1+s}(D)},$$

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where we tacitly absorbed the deterministic constant in (22) into $\mathcal{C}$. Then, the second inverse estimate in (22), (24), and the choice $\eta := h_j^{2(1-s)}$ results in

$$a(w, \phi) \leq \mathcal{C} \sqrt{\frac{\lambda_{\max}(\mathcal{A}_j)}{\check{a}}} 2 \left(\frac{h_j}{h_i}\right)^s \|\phi\|_{L^2(D)} \|w\|_{H^1(D)},$$

where we tacitly absorbed the deterministic constant in (24) into $\mathcal{C}$ as well as the constant in (22). Since

$$\frac{\|\mathcal{A}_j w\|_{L^2(D)}^2}{\lambda_{\max}(\mathcal{A}_j)} = \lambda_{\max}(\mathcal{A}_j)^{-1} \left(\sup_{\phi \in \mathcal{V}_j} \frac{a(w, \phi)}{\|\phi\|_{L^2(D)}}\right)^2,$$

we conclude with (20) and the assumption of the lemma that for every $w \in \mathcal{V}_i$

$$(T_j w, w)_{\mathcal{A}} \leq c_1 \frac{\|\mathcal{A}_j w\|_{L^2(D)}^2}{\lambda_{\max}(\mathcal{A}_j)} \leq c_1 4 \mathcal{C}^2 \left(\frac{1}{\check{a}}\right)^2 \left(\frac{h_j}{h_i}\right)^{2s} (w, w)_{\mathcal{A}}. \tag{25}$$

We argue in a similar fashion as in the second part of the proof of [25, Lemma 6.3], i.e., we conclude with the Cauchy–Schwarz inequality, (25), and with the scaling property of the smoothers in ($\mathbf{A}$) that

$$\begin{aligned}
(T_j v, T_i w)_{\mathcal{A}} &\leq (T_j v, v)_{\mathcal{A}}^{1/2} (T_j T_i w, T_i w)_{\mathcal{A}}^{1/2} \\
&\leq \mathcal{C}^2 \left(\frac{1}{\check{a}}\right)^2 (T_j v, v)_{\mathcal{A}}^{1/2} \left(\frac{h_j}{h_i}\right)^{2s} (T_i w, T_i w)_{\mathcal{A}}^{1/2} \\
&\leq \mathcal{C}^2 \eta \left(\frac{1}{\check{a}}\right)^2 \left(\frac{h_j}{h_i}\right)^{2s} (T_j v, v)_{\mathcal{A}}^{1/2} (T_i w, w)_{\mathcal{A}}^{1/2},
\end{aligned}$$

where we again absorbed deterministic constants into $\mathcal{C}$. Since $h_j/h_i \leq \gamma^{j-i}$ by assumption, the statement of the proposition follows with Lemma 2.5. $\qquad\square$

If the coefficients of $\mathcal{A}$ restricted to D_k, i.e., $(a_{ij}|_{D_k}, i, j = 1, \ldots, d)$, take values in $W^{1,\infty}(D_k)$, $k = 1, \ldots, K$, then the proof of Proposition 4.2 simplifies and enables $s = 1/2$ in Proposition 4.2, cp. [6, Remark 4.5].

Proposition 4.3 *Let the smoothers $(R_j : j = 1, \ldots, J)$ satisfy Assumption $\mathbf{B}$ with a determistic constant c_0. There exists a deterministic constant C independent of j such that inequality (10) from Assumption 1 holds with the random variable*

$$K_0 := C \frac{\hat{a}}{\check{a}}.$$

Proof. Since the assumption implies that $(R_i^{-1} v_i, v_i) \leq \lambda_{\max}(\mathcal{A}_i)/c_0 \ (v_i, v_i)$, the claim of the proposition follows with (20), (21), and (23). $\qquad\square$

4.2 Verification of Assumption 2

The proof of the regularity and approximation assumption here with explicit depends of the constants on the random operator stems from [10, Section 9.1]. Let $(H^s, \|\cdot\|_s)$, $s \in [0, 2]$, be nested Hilbert spaces, i.e., $H^s \subset H^{s'}$ for every $s > s'$, respresenting smoothness with $H^1 = \mathcal{V}$ such that the norms $\|\cdot\|_1$ and $\|\cdot\|_{\mathbb{E}(\mathcal{A})}$ are equivalent. Let $(H^{-s}, \|\cdot\|_{-s})$, $s \in (0, 2]$ denote the respective dual spaces and indentity H^0 with its dual $(H^0)^*$.

Proposition 4.4 *If there exists a random variable C_1 such that for $\mathbb{P}$-a.e. ω, for every $f \in H^{-1+s}$ the solution u of (1) takes values in H^{1+s} and satisfies the regularity estimate*

$$\|u(\omega)\|_{1+s} \le C_1(\omega)\|f\|_{-1+s}$$

and if there exists a random variable C_2 such that for $\mathbb{P}$-a.e. ω, every $v \in H^{1+s}$ is approximated by its Galerkin projection, i.e.,

$$\|(\mathrm{Id} - \mathcal{P}_{j-1}(\omega))v\|_{\mathcal{A}(\omega)} \le C_2(\omega)(\lambda_{\max}(\mathcal{A}_j(\omega)))^{-s/2}\|v\|_{1+s}$$

then, Assumption 2 holds with the random variable

$$C_s := C\hat{a}^{3-s}C_1^2 C_2^2,$$

where $C > 0$ is some deterministic constant.

Proof. We follow the proof of [10, Lemma 9.1] with the specific aim to make the dependence of the constant C_s on the operator $\mathcal{A}$ explicit. Let $v \in M_j$ be arbitrary, define $g := \mathcal{A}_j^{1-s}(\mathrm{Id} - \mathcal{P}_{j-1})v$, and let w solve $\mathcal{A}w = g$ uniquely. Note that by definition of g

$$(\mathcal{A}_j^{1-s}(\mathrm{Id} - \mathcal{P}_{j-1})v, (\mathrm{Id} - \mathcal{P}_{j-1})v) = ((\mathrm{Id} - \mathcal{P}_{j-1})v, g). \tag{26}$$

We follow the argument that proves [10, Equation (9.9)] to conclude with both assumptions of the proposition and (2) that

$$
\begin{aligned}
((\mathrm{Id} - \mathcal{P}_{j-1})v, g) &= (\mathcal{A}w, (\mathrm{Id} - \mathcal{P}_{j-1})v) \\
&\le \|v\|_{\mathcal{A}} \|(\mathrm{Id} - \mathcal{P}_{j-1})w\|_{\mathcal{A}} \\
&\le C_1 C_2 \hat{a}^{(1+s)/2}(\lambda_{\max}(\mathcal{A}_j))^{-s/2}\|v\|_{\mathcal{A}}\|g\|_{-1+s}.
\end{aligned}
\tag{27}
$$

By the definition of $\mathcal{A}_j$ and Q_j and (2), there exists a deterministic C such that for every $v \in H^1$

$$(\mathcal{A}_j Q_j v, Q_j v)^{1/2} \le C\sqrt{\hat{a}}\|v\|_1.$$

Then, by interpolation, e.g. [10, Theorem A.4], it follows that

$$(\mathcal{A}_j^{1-s}Q_j\phi, Q_j\phi)^{1/2} \le (\hat{a}C)^{1-s}\|\phi\|_{1-s} \quad \text{for every } \phi \in H^1,$$

which thereafter implies with the Cauchy–Schwarz inequality that

$$
\begin{aligned}
|(g, \phi)| = |(g, Q_j\phi)| &\le (\mathcal{A}_j^{-1+s}g, g)^{1/2}(\mathcal{A}_j^{1-s}Q_j v', Q_j v)^{1/2} \\
&\le (\sqrt{\hat{a}}C)^{1-s}(\mathcal{A}_j^{-1+s}g, g)^{1/2}\|\phi\|_{1-s}.
\end{aligned}
\tag{28}
$$

By duality it holds that $\|g\|_{-1+s} = \sup_{\phi \in H^1}(g, \phi)/\|\phi\|_{1-s}$, which implies with (28) and the definition of g that

$$\|g\|_{-1+s} \le (\hat{a}C)^{1-s}(\mathcal{A}_j^{1-s}(\mathrm{Id} - \mathcal{P}_{j-1})v, (\mathrm{Id} - \mathcal{P}_{j-1})v)^{1/2}. \tag{29}$$

Combining (26), (27), and (29) we have shown that

$$(\mathcal{A}_j^{1-s}(\mathrm{Id} - \mathcal{P}_{j-1})v, (\mathrm{Id} - \mathcal{P}_{j-1})v) \le \hat{a}^{3-s}C^{2(1-s)}C_1^2 C_2^2(\lambda_{\max}(\mathcal{A}_j))^{-s}(\mathcal{A}v, v),$$

which corresponds to the assumption [10, Equation (A.2)]. This condition is sufficient to conclude the regularity and approximation assumption with the given expression for C_s by [10, Lemma 3.2]. $\qquad\square$

Remark 4.5 *For uniformly refined FE discretizations, the assumption on the random variable C_2 can be proven by Galerkin orthogonality and an equivalence between the mesh width and the maximal eigenvalue of the stiffness matrix. In the case of the elliptic operators the regularity estimate has been derived explicitly in the coefficients of elliptic operators in Euclidean domains in [12, 23]. Respective statements for bounded, closed, smooth submanifolds in $\mathbb{R}^3$ are given in [18, Appendix C].*

5 Domain decomposition methods

We will mainly consider overlapping domain decomposition methods in the setting of Assumption (1). We denote by $\mathcal{V}_0 \subset \mathcal{V}$ a *coarse* first order FE space with meshwidth h_0 of the type introduced in Section 4.1. The first order FE space with a *fine* grid is denoted by $\widetilde{\mathcal{V}} \subset \mathcal{V}$. For a given set of overlapping subdomains $(D_j : j = 1, \ldots, J)$ of D such that $\overline{D} = \bigcup_{j=1}^{J} \overline{D}_j$. These subdomains result for example by extending a given disjoint set of subdomains by a multiple of h_0 in each spatial direction such that the union of its closures contains D. Also we assume that the boundary ∂D_j aligns with the considered mesh, $j = 1, \ldots, J$. The FE spaces $\mathcal{V}_j$, $j = 1, \ldots, J$, are subspaces of $\widetilde{\mathcal{V}}$ and are defined by

$$\mathcal{V}_j := \{v \in \widetilde{\mathcal{V}} : v(x) = 0 \ \forall x \in D \backslash D_j\}, \quad j = 1, \ldots, J$$

So, we consider the redundant space decomposition

$$\widetilde{\mathcal{V}} = \sum_{j=0}^{J} \mathcal{V}_j.$$

We consider the case that symmetric multigrid solvers from Section 4 are used as so called *subpaces solvers* $(R_j : j = 0, \ldots, J)$, which are random here. Therefore, suppose that the spaces $\mathcal{V}_j$ have nested subspaces $\mathcal{M}_{j,1} \subset \ldots \subset \mathcal{M}_{j,J'(j)} = \mathcal{V}_j$, $j = 0, \ldots, J$. Naturally, only few levels are used on the subspace $\mathcal{V}_0$, i.e. $J'(0) \approx 1$. As in Section 4, we seek for random variables K_0 and K_1 with explicit dependence on the operator $\mathcal{A}$, in order to obtain $L^q(\Omega)$-estimates for the condition numbers using additive Schwarz preconditioners and tail bounds of the random contraction number of multiplicative domain decomposition methods.

Proposition 5.1 *There exists a deterministic constant $C > 0$ that is independent of J and $J'(j)$, $j = 0, \ldots, J$, such that inequality (10) in Assumption 1 holds with the random variable*

$$K_0 := C \left(\frac{\hat{a}}{\check{a}}\right)^4 \frac{(\hat{a})^2}{(\check{a})^6} \left(1 + \left(\max_{k=1,\ldots,K} \sum_{i,j=1}^{d} \|a_{ij}\|_{W_p^s(D_k)} + \|a\|_{L^\infty(D)}\right)^4\right).$$

Proof. By (20) and [25, Lemmas 4.5 and 7.1],

$$K_0 \leq \frac{\hat{a}}{\check{a}} \frac{1}{\min_{j=0,\ldots,J} \lambda_{\min}(R_j \mathcal{A}_j)}.$$

Since the R_j's are chosen to be symmetric multigrid solvers, Propositions 4.3 and 4.2, Theorem 2.3, and (16) imply there esists a deterministic constant $c > 0$ such that for every $j = 0, \ldots, J$,

$$\lambda_{\min}(R_j \mathcal{A}_j) \geq c \frac{(\check{a})^5}{\hat{a}} \left(1 + \left(\max_{k=1,\ldots,K} \sum_{i,j=1}^{d} \|a_{ij}\|_{W_p^s(D_k)} + \|a\|_{L^\infty(D)}\right)^4\right)^{-1}.$$

This implies the assertion of the proposition. □

Proposition 5.2 *Inequality (11) holds with the deterministic number*

$$K_1 := (1 + |\{(i,j) \in \{1, \ldots, J\}^2 : D_i \cap D_j\}|)$$

Proof. The assertion will follow by [25, Lemma 4.7] after we show an estimate of the form of Assumption (**A**). By (15), it holds that $\lambda_{\max}(R_j \mathcal{A}_j) \leq 1$, $j = 0, \ldots, J$. Let $j = 0, \ldots, J$ be arbitrary. Since R_j is a symmetric multigrid solver, $R_j \mathcal{A}_j$ is symmetric and positive definite with respect to $(\cdot, \cdot)_{\mathcal{A}}$. There exists an orthonormal basis of eigenvectors of $R_j \mathcal{A}_j$ with respect to $(\cdot, \cdot)_{\mathcal{A}}$ such that $R_j \mathcal{A}_j v_i = \lambda_i v_i$. Hence, for every $v \in \mathcal{V}_j$, $v = \sum_i (v, v_i)_{\mathcal{A}} v_i$ and

$$(R_j \mathcal{A}_j v, R_j \mathcal{A}_j v)_{\mathcal{A}} = \sum_i \lambda_i^2 ((v, v_i)_{\mathcal{A}})^2 \leq \lambda_{\max}(R_j \mathcal{A}_j) \sum_i \lambda_i ((v, v_i)_{\mathcal{A}})^2 = (R_j \mathcal{A}_j v, v)_{\mathcal{A}},$$

which implies the estimate of the proposition with [25, Lemma 4.7]. □

Remark 5.3 *Nonoverlapping domain decomposition methods can also be formulated in the general framework introduced in [25, Section 3], that is recalled in Section 2, relying on Assumption 1, cp. [27, Section 2.2]. Strong convergence of PCG with deterministic iteration numbers using preconditioners for random* interface operators *of nonoverlapping domain decompositions resulting in uniform condition numbers, cp. [27, Sections 7.1 and 7.2], can also be established with Lemma 2.1 and either Theorem 3.2 or Theorem 3.4.*

6 Application to lognormal diffusion problems

The presented theory in Sections 3, 4, and 5 applies to lognormal diffusion problems. Let Z be a Gaussian random field (GRF for short) on D that takes values in Hölder spaces $C^t(\overline{D})$ such that $Z \in L^q(\Omega; C^t(\overline{D}))$ for every $q \in [1, +\infty)$ and for some $t \in (0, 1]$, for example cp. [11]. For the lognomal coefficient $a := \exp(Z)$, we consider the elliptic diffusion probem with Dirichlet boundary conditions in variational form: find $u : \Omega \to \mathcal{V}$ such that for $\mathbb{P}$-a.e. ω

$$a_\omega(u(\omega), v) = \int_D a(\omega) \nabla u(\omega) \cdot \nabla v = {}_{\mathcal{V}^*}\langle f, v \rangle_{\mathcal{V}} \quad \forall v \in \mathcal{V}, \tag{30}$$

where $\mathcal{V} = H_0^1(D)$. Well-posedness and approximation by Finite Elements is well known, cp. [11, 12, 23]. We use the FE spaces $\mathcal{V}_\ell$ from Section 4.1 with maximal meshwidth h_ℓ of $\mathcal{T}_\ell$, $\ell \in \mathbb{N}$ and remark that for each ℓ, the space $\mathcal{V}_\ell$ may have the additional structure for overlapping domain decomposition methods with multigrid subspace solvers as introduced in Section 5.

Elements of $H^{1+s}(D)$, $s \in [0, 1]$, can be approximated by functions in $\mathcal{V}_\ell$, cp. [13, Theorem 3.2.1], i.e., there exists a deterministic constant $C > 0$ such that for every $v \in H^{1+s}(D)$ there is $w_\ell \in \mathcal{V}_\ell$

$$\|v - w_\ell\|_{\mathcal{V}} \leq C h_\ell^s \|v\|_{H^{1+s}(D)}. \tag{31}$$

Note that the approximation property stated in [13, Theorem 3.2.1] can be interpolated to also hold for non-integer order Sobolev spaces. The following regularity estimate makes the dependence on the coefficient a explicit. For every $s \in [0, \min\{t, t_{-\Delta}\}) \backslash \{1/2\}$ there exists a deterministic constant $C > 0$ such that for $\mathbb{P}$-a.e. ω

$$\|u(\omega)\|_{H^{1+s}(D)} \leq C \frac{\|a(\omega)\|_{C^0(\overline{D})} \|a(\omega)\|_{C^t(\overline{D})}^2}{(\min_{x \in \overline{D}} a(\omega, x))^4} \|f\|_{H^{-1+s}}, \tag{32}$$

where $t_{-\Delta}$ is the maximal value such that the inverse of the Dirichlet Laplacian satisfies $(-\Delta)^{-1} : H^{-1+t_{-\Delta}}(D) \to \mathcal{V} \cap H^{1+t_{-\Delta}}(D)$ is bounded. For $d = 2$, the estimate (32) is due to [23, Lemma 5.2] (for $d = 3$, the reader is referred to [23, Remark 5.2(c)]). The solution u can be approximated in $\mathcal{V}_\ell$ by the FE approximation denoted by u_ℓ, in the $L^q(\Omega, \mathcal{V})$-norm, $\ell \in \mathbb{N}$.

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Specifically by Céa's lemma, (31), and (32), there exists a deterministic constant $C > 0$ that is independent of $\ell \in \mathbb{N}$ and a such that for $\mathbb{P}$-a.e. ω

$$\|u(\omega) - u_\ell(\omega)\|_{L^q(\Omega;\mathcal{V})} \leq C \left\| \frac{\|a(\omega)\|^2_{C^0(\overline{D})} \|a(\omega)\|^2_{C^t(\overline{D})}}{(\min_{x \in \overline{D}} a(\omega, x))^5} \right\|_{L^q(\Omega)} \|f\|_{H^{-1+s}(D)} \, h_\ell^s, \quad \ell \in \mathbb{N},$$

Since the embedding $C^t(\overline{D}) \subset W^{s,p}(D)$ is continuous for every $0 < s < t$ and every $p \in [1, +\infty)$, the conditions from Sections 4 and 5 are satisfied. Let

$$U_\ell^n : \Omega \to \mathcal{V}_\ell \tag{33}$$

be the result of $n \in \mathbb{N}$ iterations of an iterative algorithm introduced in the previous sections, with initial guess $U_{\ell-1}^n$, $2 \leq \ell \in \mathbb{N}$, and $U_1^n = u_1$. The iterative methods are (symmetric) multigrid (see Algorithms 4 and 3) or PCG using the additive (BPX) preconditioner (see Algorithm 1) in Section 4. In the setting of Section 5, $U_\ell^n : \Omega \to \mathcal{V}_\ell$ in (33) may result from $n \in \mathbb{N}$ iterations of PCG using the additive Schwarz preconditioner (see Algorithm 1) or a multiplicative domain decomposition method (see Algorithm 2), where symmetric multigrid (see Algorithm 3) is used as subspace solvers.

Theorem 6.1 *For $0 < s < t \leq 1$ and every $q, r \in [1, +\infty)$, there exists constant $C > 0$ such that for every number of iterations $n \in \mathbb{N}$ and $\ell \geq 2$,*

$$\|u - U_\ell^n\|_{L^q(\Omega;\mathcal{V})} \leq C(h_\ell^s + n^{-r}).$$

Proof. By the Cauchy–Schwarz inequality, $\|a\|^2_{C^0(\overline{D})} \|a\|^2_{C^t(\overline{D})} (\min_{x \in \overline{D}} a(x))^{-5} \in L^p(\Omega)$ for every $p \in [1, +\infty)$, which is one of the conditions of Theorem 3.2.

It remains to verify the needed properties of the random contraction number in the conditions of Theorem 3.2. In the framework of Assumption 1, by Theorem 2.3 the random contraction number δ satisfies for the multiplicative iteration $1/(1-\delta) \leq K_0(1+K_1)^2/(2-\nu)$. In the case of multigrid, Propositions 4.2 and 4.3 and the Cauchy–Schwarz inequality imply that $1/(1-\delta) \in L^{p'}(\Omega)$ for every $p' \in [1, +\infty)$. For overlapping domain decomposition methods, this statement is due to Propositions 5.1 and 5.2. If the additive preconditioner is applied with PCG, the random contraction number δ satisfies by Theorem 2.3, $1/(1-\delta) \leq \sqrt{K_0 K_1} + 1$. By the same argument, $1/(1-\delta) \in L^{p'}(\Omega)$ for every $p' \in [1, +\infty)$. In the setting of Assumption 2, the assumptions of Proposition 4.4 are satisfied by (31) and Galerkin orthogonality and (32). By a similar argument using Theorem 2.6, $1/(1-\delta) \in L^{p'}(\Omega)$ for every $p' \in [1, +\infty)$.

Hence, the parameter r in Theorem 3.2 may be arbitrarily large, which implies the assertion. $\square$

Corollary 6.2 *In the setting of Theorem 6.1, let U_ℓ^n result from $n \in \mathbb{N}$ iterations of PCG with a deterministic preconditioner $\widetilde{B}_j$ such that $\kappa(\widetilde{B}_j \mathbb{E}(\mathcal{A}_j))$ is bounded uniformly in j. Then, the strong convergence estimate of Theorem 6.1 also holds.*

Proof. By Lemma 2.1, for $\mathbb{P}$-a.e. ω, $\kappa(\widetilde{B}_j \mathcal{A}_j(\omega)) \leq \hat{a}(\omega)/\check{a}(\omega) \kappa(\widetilde{B}_j \mathbb{E}(\mathcal{A}_j))$. Since $\hat{a}/\check{a} \in L^{q'}(\Omega)$ for every $q' \in [1, +\infty)$, the claim follows as in the proof of Theorem 6.1. $\square$

Corollary 6.3 *In the setting of Theorem 6.1, for every $\varepsilon > 0$ there exists a constant $C_{q,\varepsilon,s} > 0$ that is independent of h_ℓ such that for every $\ell \geq 2$*

$$\|u - U_\ell^n\|_{L^q(\Omega;\mathcal{V})} \leq C_{q,\varepsilon,s} \, h_\ell^s.$$

with a deterministic number of iterations given by $n = \lceil c h_\ell^{-\varepsilon} \rceil$ for a deterministic constant $c > 0$. The cost for one sample of U_ℓ^n is

$$\mathcal{O}(h_\ell^{-d-\varepsilon})$$

with deterministic constants that are independent of h_ℓ, $\ell \geq 0$.

Proof. The cost of one iteration is $\mathcal{O}(h_\ell^{-d})$, since the matrix vector product has cost $\mathcal{O}(h_\ell^{-d})$ for the sparse stiffness matrices that result taking the nodal basis of $\mathcal{V}_\ell$. The error contributions in the estimate of Theorem 6.1 are equilibrated for this choice of iteration number, since r in Theorem 6.1 can be chosen arbitrarily large, i.e., $r = s/\varepsilon$ is admissible. $\qquad\square$

6.1 Mean field approximation with MLMC

Let $G : \Omega \times \mathcal{V} \to \mathcal{W}$ be a random operator, where $\mathcal{W}$ is a Banach space, that is Lipschitz continuous, i.e., there exists a positive random variable $\mathcal{C}$ such that for $\mathbb{P}$-a.e. ω

$$\|G(\omega, v) - G(\omega, w)\|_{\mathcal{W}} \leq \mathcal{C}(\omega)\|v - w\|_{\mathcal{V}} \quad \forall v, w \in \mathcal{V}.$$

Let us assume that $\mathcal{C} \in L^{q'}(\Omega)$ for some $q' \in (2, +\infty)$. For the computation of the mean field $\mathbb{E}(G(u))$, MLMC methods are well established, cp. [12, 15]. For a maximum number of levels $L \in \mathbb{N}$, define

$$E_L^{\mathrm{ML}}(G(U_L^n)) := \sum_{\ell=1}^{L} E_{M_\ell}(G(U_\ell^n) - G(U_{\ell-1}^n)),$$

where $(E_{M_\ell} : \ell = 1, \ldots, L)$ are Monte Carlo estimators that are mutually independent. We used the convention that $U_0^n := 0$. For the effectiveness of MLMC, bounds on the $L^2(\Omega; \mathcal{V})$-norm of the difference between two consecutive discretization levels have to be known. Specifically, by Theorem 6.1 and Corollary 6.3, the Hölder inequality, and the triangle inequality

$$\|G(U_\ell^{n_\ell}) - G(U_{\ell-1}^{n_{\ell-1}})\|_{L^2(\Omega;\mathcal{W})} \leq \|\mathcal{C}\|_{L^{q'}(\Omega)}\|U_\ell^{n_\ell} - U_{\ell-1}^{n_{\ell-1}}\|_{L^{2q'/(q'-2)}(\Omega;\mathcal{V})} = \mathcal{O}(h_{\ell-1}^s),$$

where $n_\ell = \lceil ch_\ell^{-\varepsilon} \rceil$ for a deterministic constant $c > 0$ and a small deterministic number $\varepsilon > 0$. Let us assume that the stiffness matrix on every level $\ell \in \mathbb{N}$ can be assembled in computational work $\mathcal{O}(h_\ell^{-d} \log(h_\ell^{-1}))$. Possible techniques to achieve this are in certain cases the circulant embedding or FFT, cp. [12, Section 5]. See also [22] for approximation techniques of the so called *Karhunen–Loève* expansion. Also, suppose that the computational cost of the evaluation of $G(v)$ is $\mathcal{O}(h_\ell^{-d})$ for every $v \in \mathcal{V}_\ell$. For example circulant embedding or FFT are applicable if $\log(a)$, which is a Gaussian random field, has stationary covariance kernel and is the restriction of such a Gaussian random field defined on a product domain containing D. Then, by the complexity theorem [12, Theorem 4.1] (see also [15]) and Corollary 6.3, there exist sample numbers $(M_\ell, \ell = 1, \ldots, L)$ such that an error threshold $0 < \mathrm{TOL} < e^{-1}$, i.e.,

$$\|\mathbb{E}(G(u) - E_L^{\mathrm{ML}}(G(U_L^n))\|_{L^2(\Omega;\mathcal{V})} = \mathcal{O}(\mathrm{TOL}),$$

can be achieved with computational cost

$$\mathrm{work}_L = \begin{cases} \mathcal{O}(\mathrm{TOL}^{-2}) & \text{if } 2s > d + \varepsilon, \\ \mathcal{O}(\mathrm{TOL}^{-(d+\varepsilon)/s}) & \text{if } 2s < d + \varepsilon, \end{cases} \tag{34}$$

where $d = 1, 2, 3$ is the dimension of the domain D. Note that the case $2s = d + \varepsilon$ is not included, since here s satisfies a strict inequality, which renders this case unimportant. The estimate on the computational work of MLMC for lognormal diffusion coefficients in (34) are deterministic and do not have assumptions on the complexity of the PDE solver. The computational work of MLMC FE for lognormal diffusions is for $d = 2, 3$ essentially optimal, since $\varepsilon > 0$ may be chosen arbitrarily small.

Remark 6.4 *The established theory is also applicable for deterministic preconditioners that do not imply uniform condition numbers. A possible class of examples are so called "algebraic*

multigrid" (AMG for short) preconditioners for which a multilevel convergence theory does not seem to be available, cp. [26, Section 5]. However, if such a deterministic AMG preconditioner could be tuned to $\mathbb{E}(\mathcal{A}_j)$, the respective iterative method also converges strongly as in Theorem 6.1 by a similar argument applying Theorem 3.4 and Corollary 6.2. Note that in the example (30) (see also Example 1), $\mathbb{E}(\mathcal{A})$ is the Dirichlet Laplacean multiplied by $\mathbb{E}(\exp(Z))$, which is constant with respect to $x \in D$ due to the assumed stationarity of the GRF Z.

7   Conclusions and extensions

In the study of random operator equations with non-uniform input a rigorous framework has been established to verify strong convergence of iterative solvers. In the case of lognormal random input, essentially optimal, deterministic complexity bounds are implied. This offers an alternative to direct solvers for this type of problems, whose performance may decline in higher dimensions, e.g., $d = 3$. Assumptions on the computational cost of PDE solvers are also common in the context of multilevel quasi-Monte Carlo methods (MLQMC for short) for PDE problems, cp. [21, 20, 14, 19]. Applicability of iterative solvers for MLQMC will be analyzed in a forthcoming work.

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